

Proof of Concurrent C Code Via Automatic Code Transformation

Motivation and Goals

Frama-C and WP allow deductive proof of sequential C code. To deal with concurrent code, we propose to transform it into sequential code.

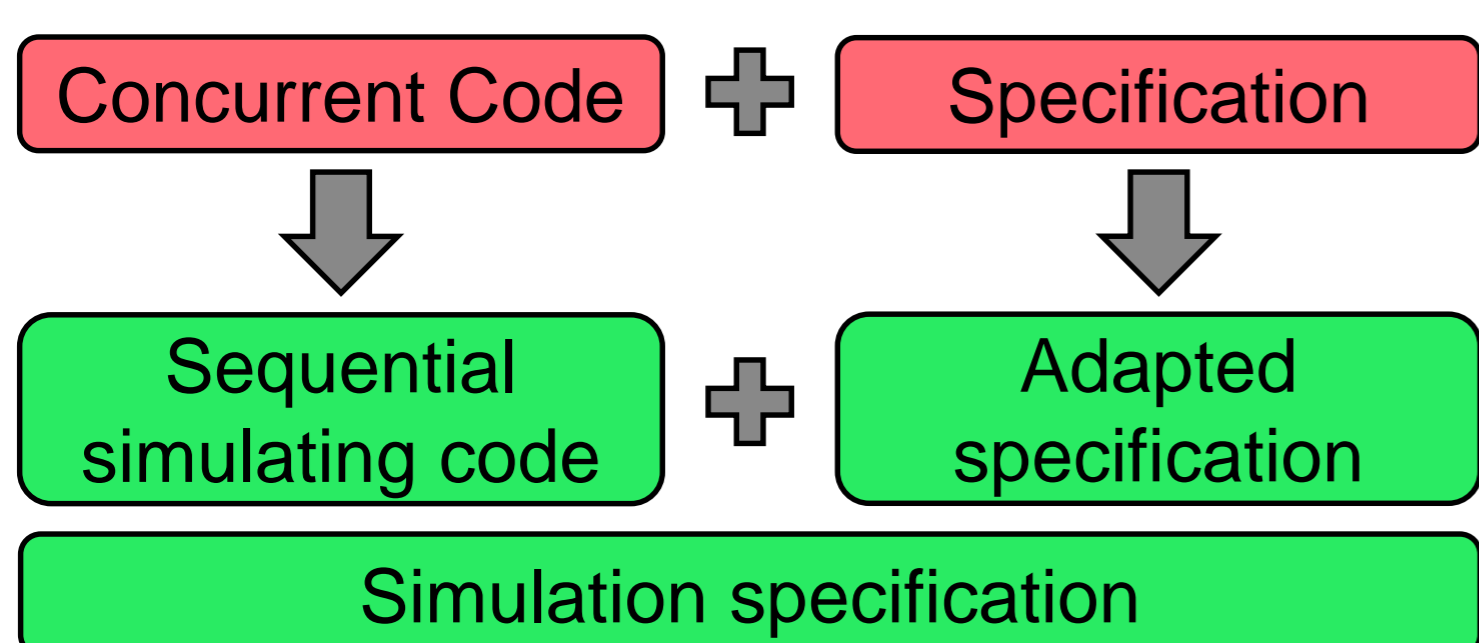
Needs :

- Automatic proof of generated code
- Automatic code transformation
- Correctness of the transformation

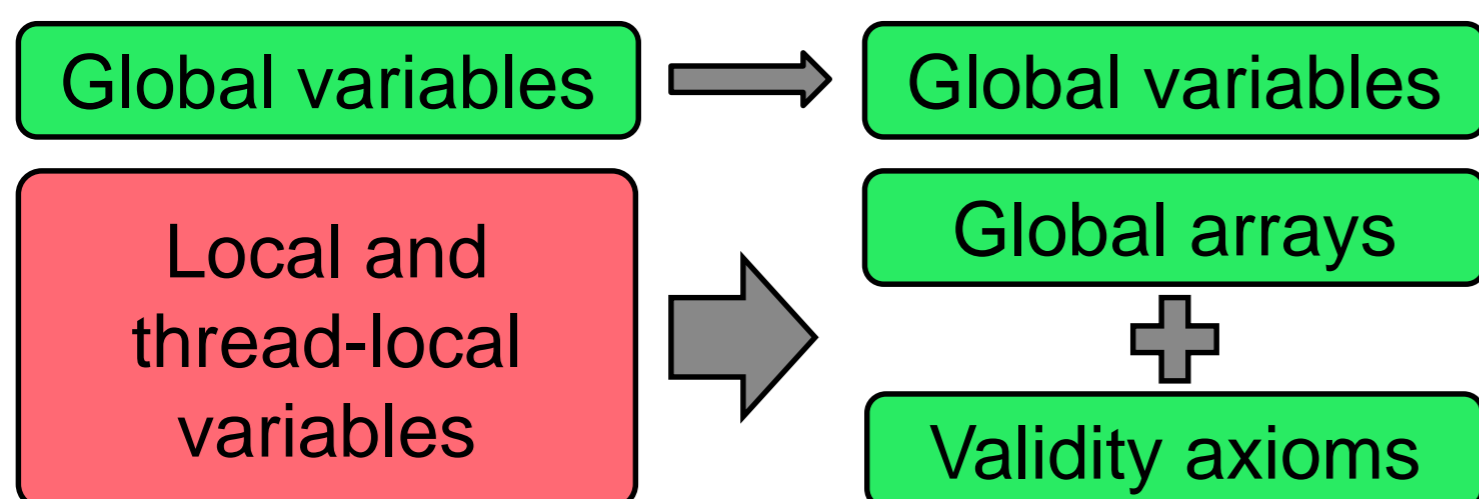
Hypotheses :

- Interleaving semantics
- No dynamic thread spawning
- User-specified atomic sections

Overview of the Code Transformation



Variables



Global Approach

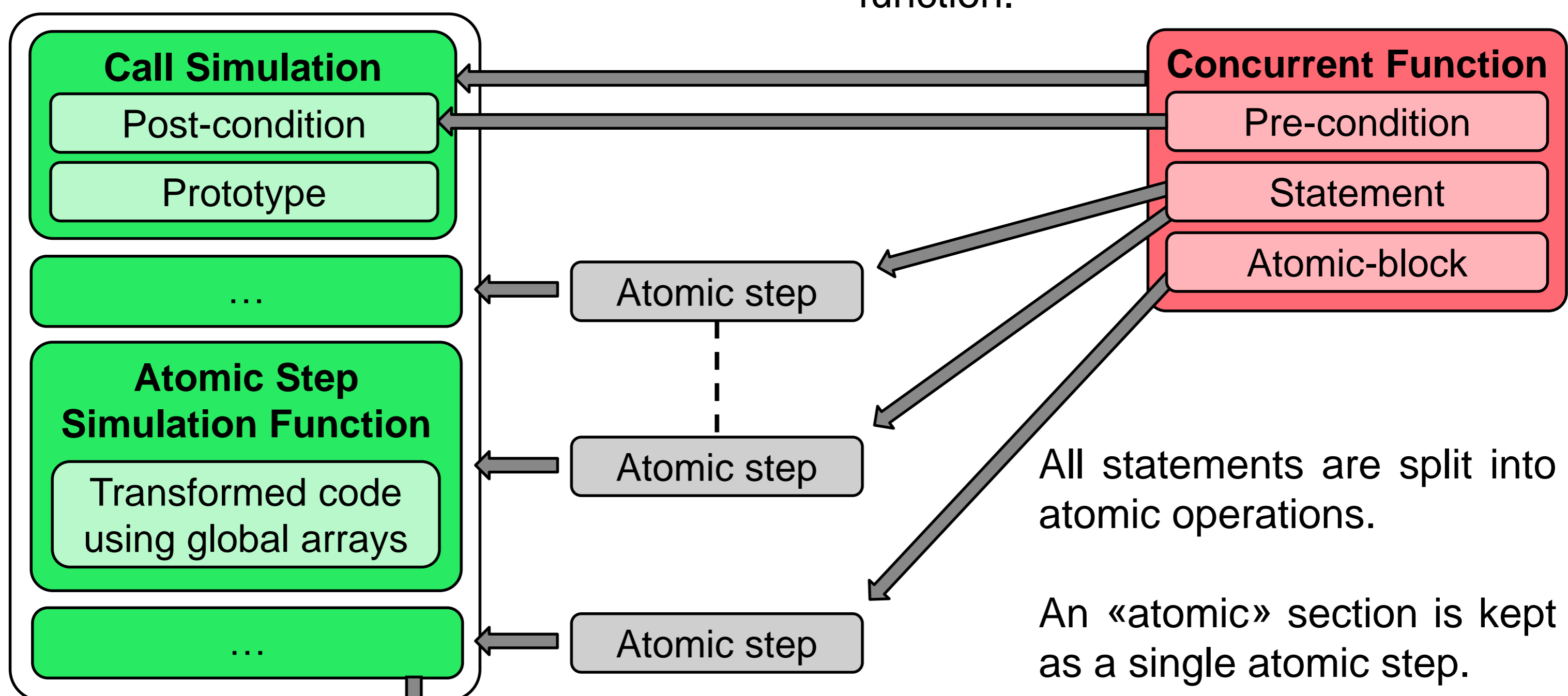
We transform the original code into a sequential simulating one. The code we get can be then proved using WP :

- Automatic proof of generated code

Functions and statements

A concurrent function call is modeled by a specified C function prototype.

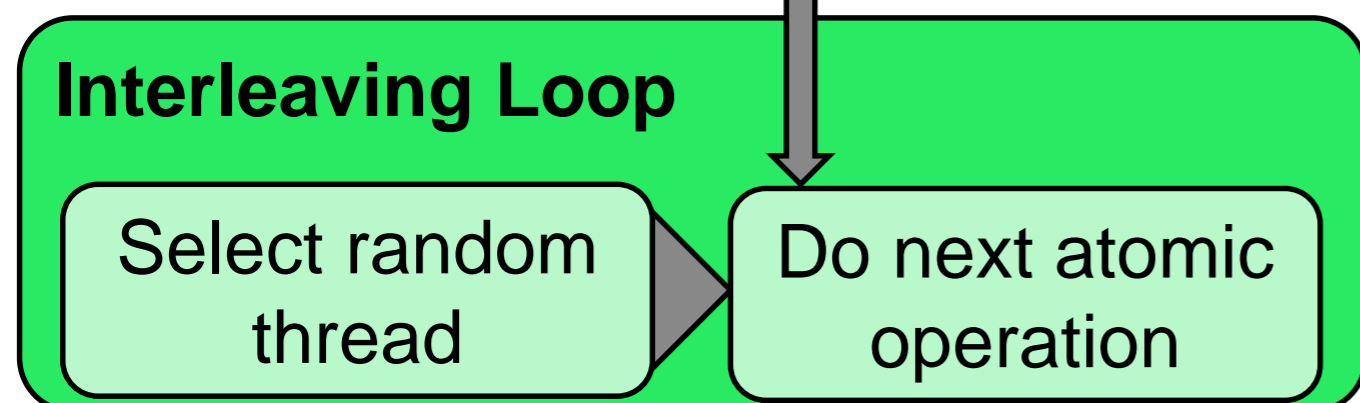
Every atomic step is simulated by a function.



All statements are split into atomic operations.

An «atomic» section is kept as a single atomic step.

Interleavings



Atomic operations are interleaved using an infinite loop.

Specifications

We state simulation invariants and global invariants as pre/post conditions of simulating functions, and as a loop invariant of the interleaving loop :

- Automatic transformation
- Code
- Specification

Ongoing and Future Work

Ongoing Work :

- Correctness of the transformation
- Code (with Coq Proof Assistant)
- Specification

Future Work :

- Add concurrency primitives to ACSL
- Support weaker semantics

Reference : A. Blanchard, N. Kosmatov, M. Lemerre, and F. Loulergue. A Case Study on Formal Verification of the Anaxagoras Hypervisor Paging System with Frama-C. FMICS 2015.

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